Problem 1: 3.2.3 (also see details in assignment document]
$\sim_{\text {avg }}(n)$.
headquarters fin wants to determine the highest floor of its $n$-story functionality. The firm a fall with no impact on the gadget's an algorithm in the best efficiency class you can to experiment with. Design
$A$ $\qquad$ ....

Let N be the total number of floors, and F the number of the lowest floor on which a gadget fails when dropped from there.
Assume that due to weight, volatility, or some other factor, there is a high cost ( $C$ ) for each floor that a gadget must be carried up. Also a cost ( $T$ ) for each test to determine whether the drop caused the gadget to fail after each drop.

First, give big-Theta worst-case running times in terms of $\mathrm{N}, \mathrm{C}$, and T for the obvious algorithm that tries every floor in succession (that algorithm only requires one gadget). Then design and analyze the most efficient algorithm you can.

Can you think of any flaws in this general approach to testing?
Problem 2: 3.2.4
4. Determine the number of character comparisons that will be made by the
brutc-force alyurithm in searching for the pattern GANDHI in the text
THERE_IS_MORE_TO_LIFE_THAN_INCREASTNC_ITS_SPEED
(Assume that the length of the text-it is 47 characters long-is known before
the search starts.)

## Problem 3: 3.2.6

6. Give an example of a text of length $n$ and a pattern of length $m$ that constitutes a worst-case input for the brute-force string-matching algorithm. Exactly how many character comparisens will be madc for such input?

- ...

Problem 4: 3.3.4 [3.3.3]
3. a. There are several alternative ways to define a distance between two points $P_{1}=\left(x_{1}, y_{1}\right)$ and $P_{2}=\left(x_{2}, y_{2}\right)$ in the Cartesian plane. In particular, the so-called Manhattan distance is defined as

$$
d_{M}\left(P_{1}, P_{2}\right)=\left|x_{1}-x_{2}\right|+\left|y_{1}-y_{2}\right| .
$$

Prove that $d_{M}$ satisfies the following axioms which every distance function must satisfy:
i. $d_{M}\left(P_{1}, P_{2}\right) \geq 0$ for any two points $P_{1}$ and $P_{2}$, and $d_{M}\left(P_{1}, P_{2}\right)=0$ if and only if $P_{1}=P_{2}$;
ii. $d_{M}\left(P_{1}, P_{2}\right)=d_{M}\left(P_{2}, P_{1}\right)$;
iii. $d_{M}\left(P_{1}, P_{2}\right) \leq d_{M}\left(P_{1}, P_{3}\right)+d_{M}\left(P_{3}, P_{2}\right)$ for any $P_{1}, P_{2}$, and $P_{3}$.
b. Sketch all the points in the $x, y$ coordinate plane whose Manhattan distance to the origin $(0,0)$ is equal to 1 . Do the same for the Euclidean distance.
c. True or false: A solution to the closest-pair problem does not depend on which of the two metrics- $d_{E}$ (Euclidean) or $d_{M}$ (Manhattan)-is used?

Clarifications: (b) You can sketch or simply list them.
(c) By "a solution", they mean "an algorithm to solve the problem".
5. The closest-pair problem can be posed on the $k$-dimensional space in which the Euclidean distance between two points $P^{\prime}=\left(x_{1}^{\prime}, \ldots, x_{k}^{\prime}\right)$ and $P^{\prime \prime}=\left(x_{1}^{\prime \prime}, \ldots, x_{k}^{\prime \prime}\right)$ is defined as

$$
d\left(P^{\prime}, P^{\prime \prime}\right)=\sqrt{\sum_{s=1}^{k}\left(x_{s}^{\prime}-x_{s}^{\prime \prime}\right)^{2}}
$$

What will be the efficiency class of the brute-force algorithm for the $k$ dimensional closest-pair problem?

## Problem 6: 3.3.10 [ 3.3.8]

8. What modification needs to be made in the brute-force algorithm for the convex-hull problem to handle more than two points on the same straight line?

## Problem 7: 3.4.1 (traveling salesman)

1. a. Assuming that each tour can be generated in constant time, what will be the efficiency class of the exhaustive-search algorithm outlined in the text for the traveling salesman problem?
b. If this algorithm is programmed on a computer that makes 1 billion additions per second, estimate the maximum number of cities for which the problem can be solved in
i. one hour.
ii. 24 -hours.
iii. one year.
iv. one century.

## Problem 8: 3.4.5

5. Give an example of the assignment problem whose optimal solution does not include the smallest element of its cost matrix.

Problem 9: 3.4.6
6. Consider the partition problem: given $n$ positive integers, partition them into two disjoint subsets with the same sum of their elements. (Of course, the problem does not always have a solution.) Design an exhaustive-search algorithm for this problem. Try to minimize the number of subsets the algorithm needs to generate.

