CSSE232 Computer Architecture I

Datapath

Class Status

- Reading
 - Sections 4.1-3

- Project
 - Project group milestone assigned
 - Indicate who you want to work with
 - Indicate who you don't want to work with
 - Due next Friday (before exam)

Review so far

- Performance
- Instruction Sets (ISAs)
- MIPS assembly
 - Register convention
 - Procedure calls
 - Alignment
 - Exceptions

Outline

- Today, we begin hardware
 - Datapath

- Stages of execution
- CPU overview
- Building a datapath

Introduction

- CPU performance factors
 - Instruction count
 - Determined by ISA and compiler
 - CPI and Cycle time
 - Determined by CPU hardware
- We will examine three MIPS datapaths
 - A simplified version
 - An improved split design
 - A more realistic pipelined version

Different Subsets of Instructions

- Memory reference: 1w, sw
- Arithmetic/logical: add, sub, and, or, slt
- Control transfer: beq, j

- Instructions executed in multiple steps
 - Some steps are common to all instructions
 - Some steps shared between subsets
 - Some steps unique to single instruction

Five Stages of Instruction Execution

- 1. Instruction fetch
- Instruction decode
- 3. Execute
- 4. Memory access
- 5. Write back

We will formalize hardware for these stages after Winter break. Until then, we'll build parts as needed.

Instruction Fetch

- Fetch the instruction
 - based on value stored in PC
- Prepare for next instruction
 - Increment the PC

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Instruction Decode

Read registers from register file

Calculate branch address



Execute

- Depending on instruction class
- Use ALU to calculate
 - Arithmetic result (R-type)
 - Memory address (lw,sw)
 - Branch compare (beq, bne)



PC ← target address (faster branches)

Memory

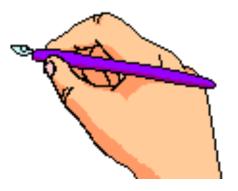
- Store data into memory (sw)
- Read data from memory (lw)

- Other things happen here too
 - Save results from R-type instructions

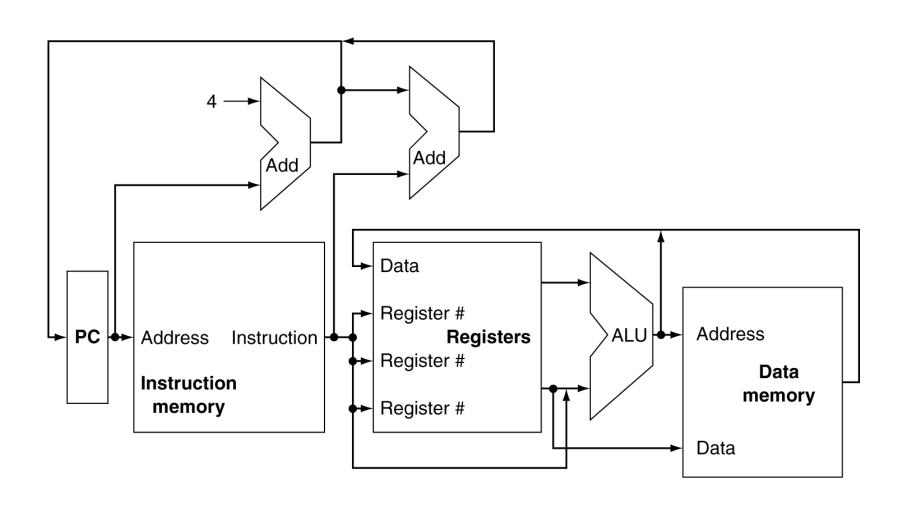


Write Back

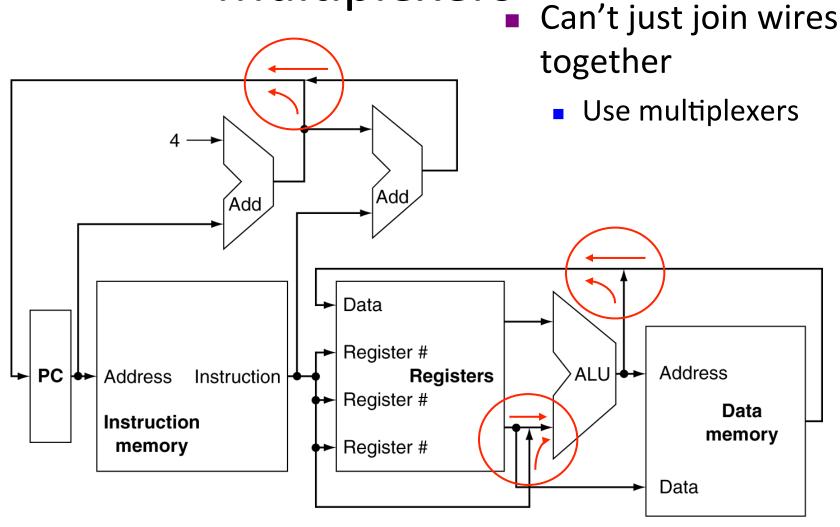
- Store the data from memory into registers (lw)
 - Similar to saving R-type results



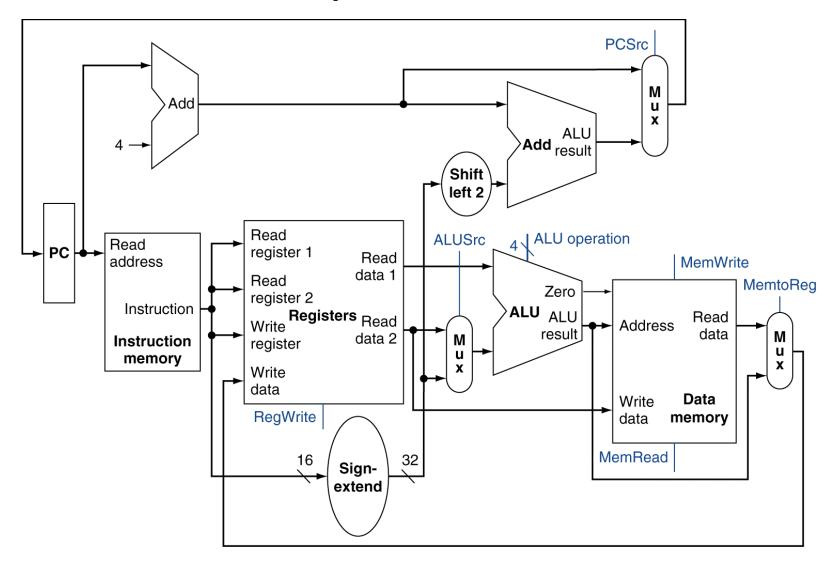
Datapath Overview

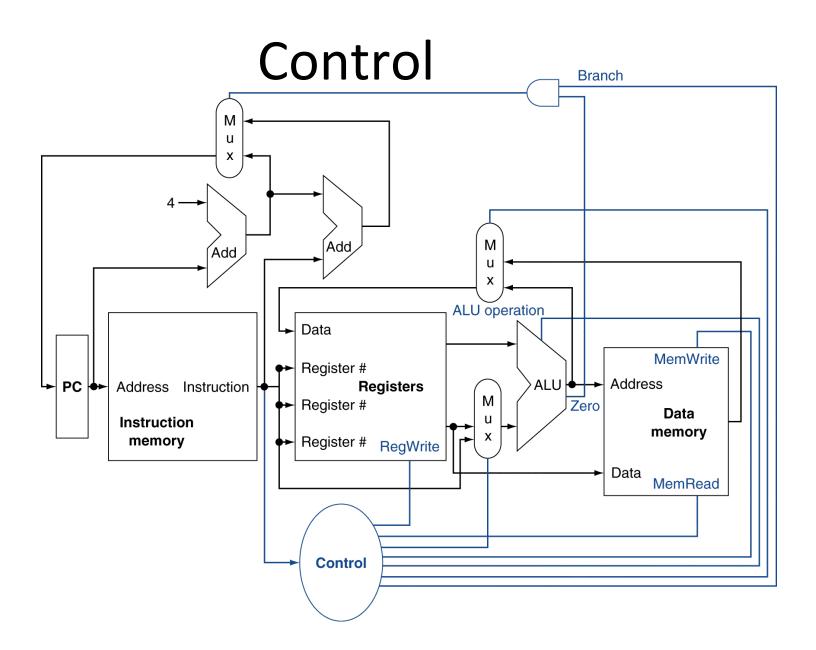


Multiplexers



Multiplexers

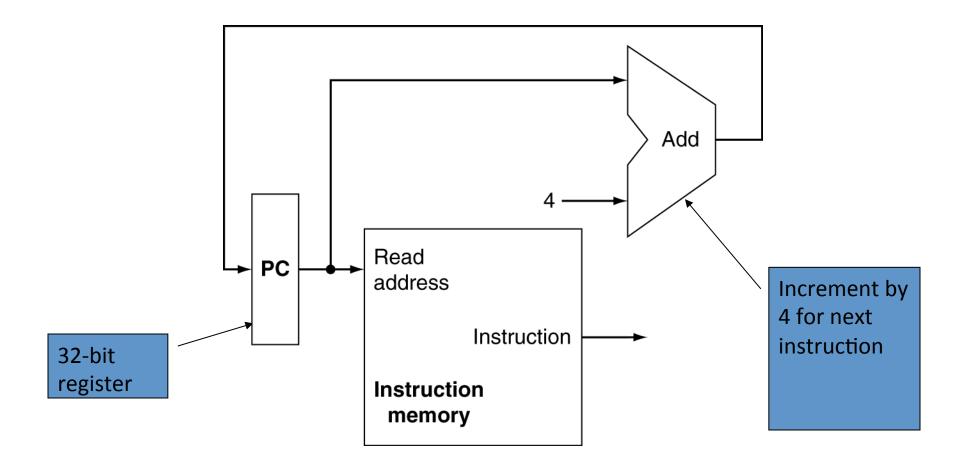




Building a Datapath

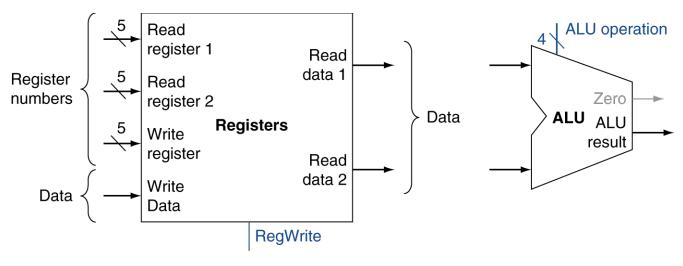
- Datapath
 - Elements that process data and addresses in the CPU
 - Registers, ALUs, mux's, memories, ...
- We will build a MIPS datapath incrementally
 - Refining the overview design
 - Consider how you will build your datapath!

Instruction Fetch



R-Format Instructions

- Read two register operands
- Perform arithmetic/logical operation
- Write register result

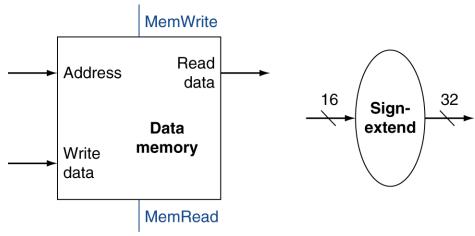


a. Registers

b. ALU

Load/Store Instructions

- Read register operands
- Calculate address using 16-bit offset
 - Use ALU, but sign-extend offset
- Load: Read memory and update register
- Store: Write register value to memory



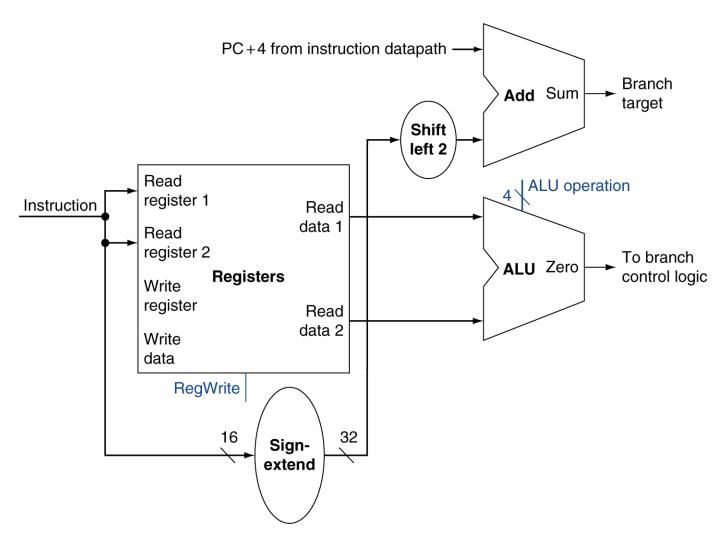
a. Data memory unit

b. Sign extension unit

Branch Instructions

- Read register operands
- Compare operands
 - Use ALU, subtract and check Zero output
- Calculate target address
 - Sign-extend displacement
 - Shift left 2 places (word displacement)
 - Add to PC + 4
 - Already calculated by instruction fetch

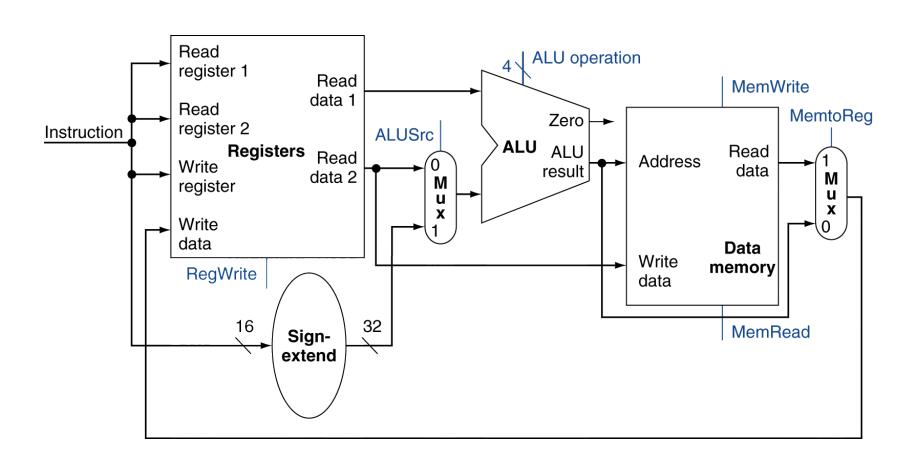
Branch Instructions



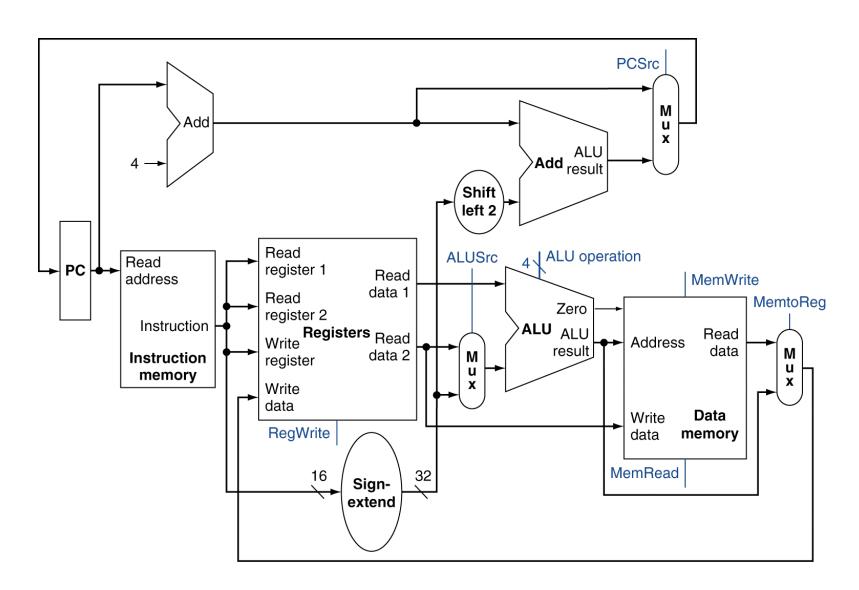
Composing the Elements

- First datapath does one instruction in one clock cycle
 - Each datapath element can only do one function at a time
 - Hence, we need separate instruction and data memories
- Use multiplexers where alternate data sources are used for different instructions

R-Type/Load/Store Datapath



Full Datapath



Performance Issues

- Longest delay determines clock period
 - Critical path: load instruction
 - InstMem→ RegFile→ ALU→ DataMem→ RegFile
- Not feasible to vary period for different instructions
- Violates design principle
 - Making the common case fast
- We will improve performance in next version

Review and Questions

- Stages of execution
- CPU overview
- Building a datapath

Different Subsets of Instructions

- Memory reference: 1w, sw
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- Control transfer: beq

Two programs...

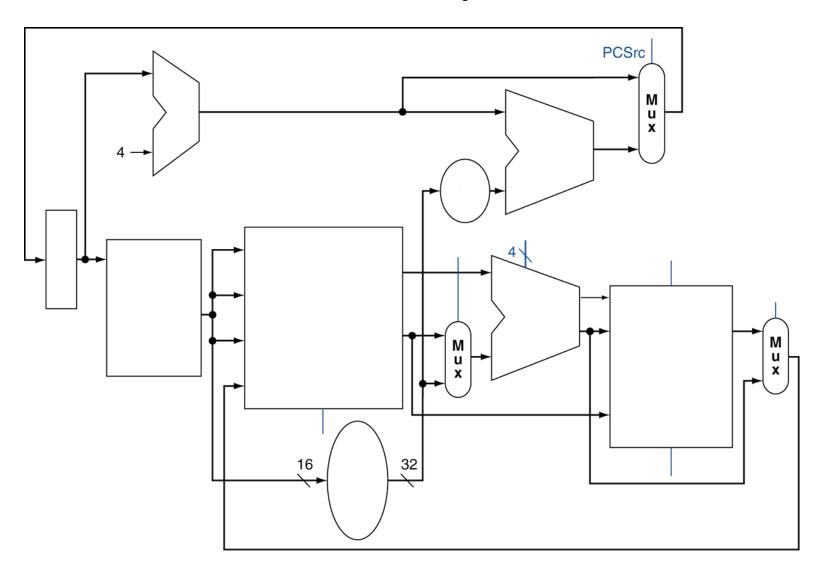
Program 1

- lw \$t0 4(\$t1)
- add \$t0 \$t0 \$t2
- beq \$t0 \$zero exit
 beq \$s1 \$t2 loop

Program 2

- sub \$s1 \$s1 \$s2
- sw \$s1 0(\$t0)

Blank Datapath



Datapath Exercise

- Trace instructions through datapath
 - Label used datapath components
 - Trace path through datapath

- Write your name at the top
- And the program you used