

Analysis of Algorithms

Functions in Increasing Order of Growth Rate

- Constant c
- Logarithmic $\log N$
- Log-squared $\log^2 N$
- Linear N
- $N \log N$
- Quadratic N^2
- Cubic N^3
- Exponential 2^N

Pigeonholing

- Realistically, we have eight choices
- Pretty good odds

Algorithms vs. Programs

- We analyze algorithms rather than programs
- Ignores poor implementation
- Focuses on the big picture

What to Analyze?

- Determine statements that contribute to the major work being done by the algorithm.
- Determine the number of times they get executed.

Worst, best, average case

- We **ALWAYS** perform an analysis for the general case of n .
- *Best case*: For input of size n , what is the **best** possible running time.
- *Worst case*: For input of size n , what is the **worst** possible running time.
- *Average case*: For input of size n , what is the **average** running time.

Linear Search

```
public static int linearSearch(int[] a, int e){
    for (int i = 0; i < a.length; i++){
        if (a[i] == e) return i;
    }
    return -1;
}
```

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Best Case Analysis of Linear Search

- Size of array is of length n .
- In **best** case, the element we are looking for is in the first position of the array.
- In this case, we have one comparison.
- $O(1)$

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Worst Case Analysis of Linear Search

- Size of array is of length n .
- In **worst** case, the element we are looking for is in the last position of the array or not located in the array
- In these cases, we have to look at all elements of the array, giving n comparison.
- $O(n)$

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Average Case Analysis of Linear Search

- Chances of looking for 1st element in array:
 $1/n$
- Same for all other elements
- Number of elements to compare:
 - 1st element: 1
 - 2nd element: 2
 - n^{th} element: n

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Average Case Analysis of Linear Search

- Sum of all cases: $1/n * 1 + 1/n * 2 + \dots + 1/n * n$
- Factor out $1/n$: $1/n * (1 + 2 + \dots + n)$
- Change notation: $1/n * \sum_{i=0}^n i$
- By induction, you can show that:

$$\sum_{i=0}^n i = n * (n+1) / 2$$
- Dividing by n : $(n+1) / 2$
- $O(n)$

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EZ Analysis

- Implement algorithm
- Count the number of times key statements get executed
- Have the computer print the best, worst and average cases.