

# CSSE 220 Day 23

Exam Review

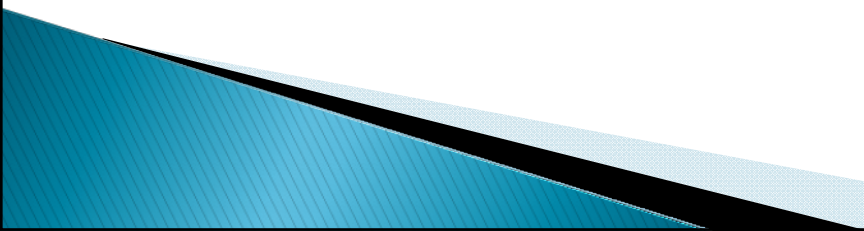
Minesweeper mine placement

Hardy Efficiency

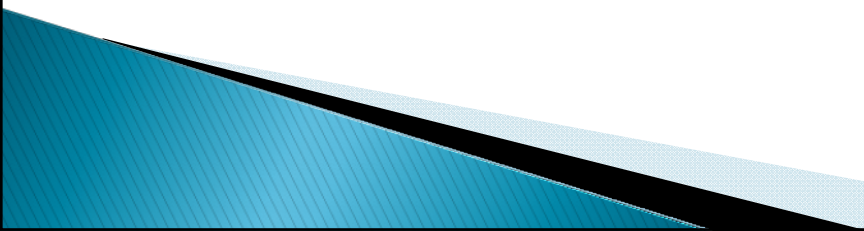
# CSSE 220 Day 23

- ▶ **Reminder: Exam #2 is this Thursday**
  - In order to reduce time pressure, you optionally may take the non-programming part 7:10–7:50 AM.
  - You may bring one piece of paper with notes for the first part.
  - Same resources as last time for the programming part.
- ▶ **Markov Milestone 2 due Friday 5 PM**
- ▶ **Begin thinking about Spell-check program**
- ▶ **You can still do the Mini-project partner surveys this morning**
- ▶ **Blood Drive today and tomorrow – Union**

# Today's Agenda

- ▶ Answers to your questions in preparation for the exam
  - ▶ Some (not-so stupid) Minesweeper tricks.
  - ▶ A look at my Hardy solution
  - ▶ Empirical analysis of an algorithm.
  - ▶ More on Linked Lists?
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# Answers to your questions

- ▶ Abstract Data Types and Data Structures
  - ▶ Collections and Lists
  - ▶ Markov
  - ▶ Exam
  - ▶ Material you have read
  - ▶ Anything else
- 

# Minesweeper tricks

- ▶ Picking random locations for mines
- ▶ Counting neighboring mines

# A Hardy Algorithm

- ▶ total =  $a^3 + b^3$ .
- ▶ One way to move through a and b loops:

a ↓ b →	1	2	3	4	5	6
0						
1						
2						
3						
4						
5						
6						

# A Hardy Algorithm

- ▶  $\text{total} = a^3 + b^3$ .
- ▶ One way to move through  $a$  and  $b$  loops:

$a \downarrow b \rightarrow$	1	2	3	4	5	6
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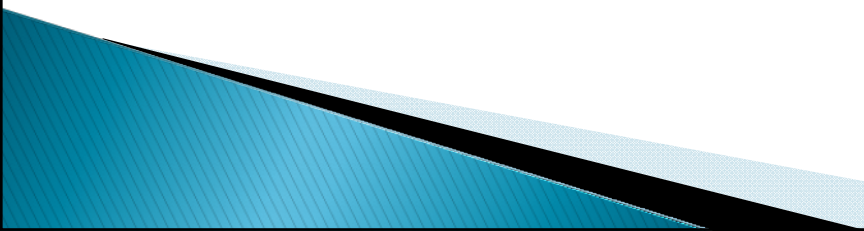
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1	■	■				
2	■	■	■			
3	■					
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5						
6						

# Hardy Algorithm basic idea

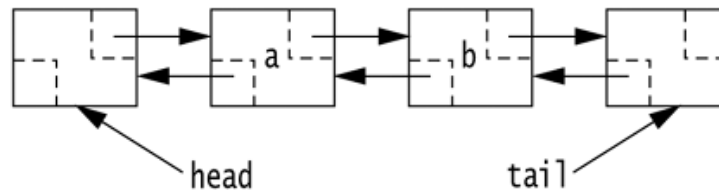
- ▶ Go through the values of  $a$  and  $b$  in the order just described
- ▶ When we calculate each total
  - Look in table if we have seen that total before
  - If not, record its triple:  $(a, b, \text{total})$  in table.
  - If so, record in the duplicates table
- ▶ When we get  $N$  items in the duplicates table
  - They may not be the  $N$  smallest. Sort them
  - See if we can find any others with sums smaller than the max of those  $N$ .
    - If, so, they will all have a  $b$  that is less than the cube root of this max. Find all of those and add to duplicates table.
- ▶ Sort again and pick out the  $N$ th one.

# Hardy Code

- ▶ Look at it together
  - ▶ Ask questions about anything you don't understand.
  - ▶ I'll ask you questions.
  - ▶ We'll add some timing computations.
  - ▶ Try to figure out a big-Oh estimate.
  - ▶ Then see how much of a speed-up we get by using a faster data structure
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# Doubly-linked list

- ▶ Each node has two pointers, **prev** and **next**.
- ▶ There is one other new node, **tail**, whose **prev** pointer points to the node containing the last element of the list.
- ▶ This makes `remove()` easier to write
  - and it also makes an efficient `ListIterator` possible.



**figure 17.15**  
A doubly linked list