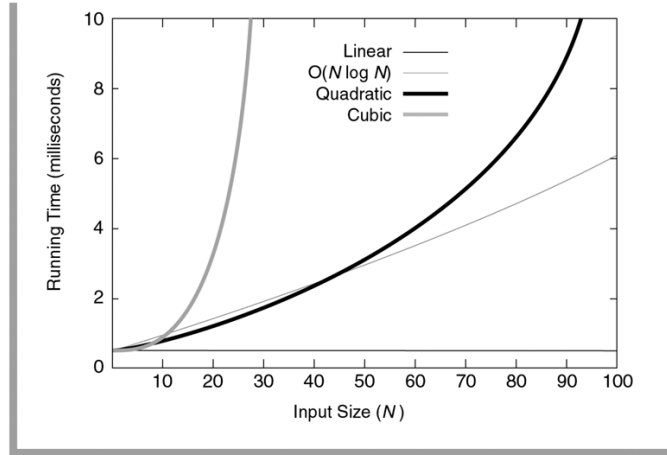
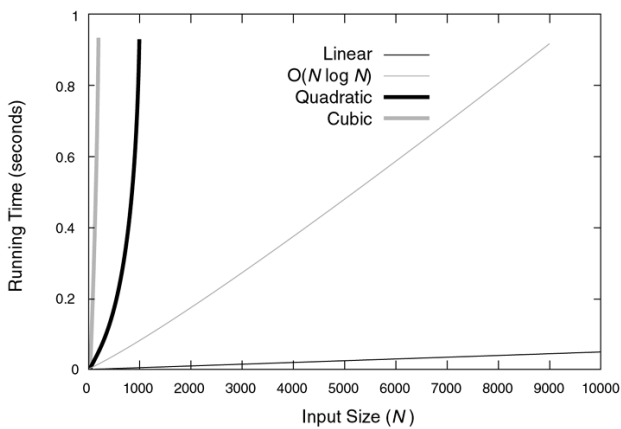


**figure 5.1**  
Running times for small inputs



Copyright © 2010 Pearson Education, publishing as Addison-Wesley. All rights reserved

1-1



**figure 5.2**  
Running times for moderate inputs

Copyright © 2010 Pearson Education, publishing as Addison-Wesley. All rights reserved

1-2

Function	Name
$c$	Constant
$\log N$	Logarithmic
$\log^2 N$	Log-squared
$N$	Linear
$N \log N$	$N \log N$
$N^2$	Quadratic
$N^3$	Cubic
$2^N$	Exponential

**figure 5.3**

Functions in order of increasing growth rate

## Basic Principles of Analysis of Algorithms

- Determine which statements or expressions to count.
- For any  $n$ , one may determine:
  - Best case
  - Average case
  - Worst case
- If the best and worst case are in the same complexity class, so is the average case.
- Average case analysis is typically the hardest. It requires a probabilistic analysis.
- We are interested in the worst case behavior, when it comes to process control.
- We are interested in the average case behavior, when it comes to minimizing hardware cost for software that is run many times.
- We might perform a best case analysis, if we want to determine the average case and suspect that the best and worst case are the same.

# Linear Search

```
public static int linearSearch(int[] a, int e){
    for (int i = 0; i < a.length; i++){
        if (a[i] == e) return i;
    }
    return -1;
}
```

Copyright : Michael Wollowski

1-5

## Best Case Analysis of Linear Search

- Size of array is of length  $n$ .
- In **best** case, the element we are looking for is in the first position of the array.
- In this case, we have one comparison.
- $O(1)$

Copyright : Michael Wollowski

1-6

## Worst Case Analysis of Linear Search

- Size of array is of length  $n$ .
- In **worst** case, the element we are looking for is in the last position of the array or not located in the array
- In these cases, we have to look at all elements of the array, giving  $n$  comparison.
- $O(n)$

Copyright : Michael Wollowski

1-7

## Average Case Analysis of Linear Search

- Chances of looking for 1<sup>st</sup> element in array:  
 $1/n$
- Same for all other elements
- Number of elements to compare:
  - 1<sup>st</sup> element: 1
  - 2<sup>nd</sup> element: 2
  - $n^{\text{th}}$  element:  $n$

Copyright : Michael Wollowski

## Average Case Analysis of Linear Search

- Sum of all cases:  $1/n*1 + 1/n*2 + \dots + 1/n*n$
- Factor out  $1/n$ :  $1/n*(1 + 2 + \dots + n)$
- Change notation:  $1/n * \sum_{i=0}^n i$
- By induction, you can show that:  
$$\sum_{i=0}^n i = n*(n+1)/2$$
- Dividing by  $n$ :  $(n+1)/2$
- $O(n)$