



figure 12.1

A standard coding scheme

| Character | Code | Frequency | Total Bits |
|--------------|------|-----------|------------|
| a | 000 | 10 | 30 |
| e | 001 | 15 | 45 |
| i | 010 | 12 | 36 |
| s | 011 | 3 | 9 |
| t | 100 | 4 | 12 |
| <i>sp</i> | 101 | 13 | 39 |
| <i>nl</i> | 110 | 1 | 3 |
| Total | | | 174 |

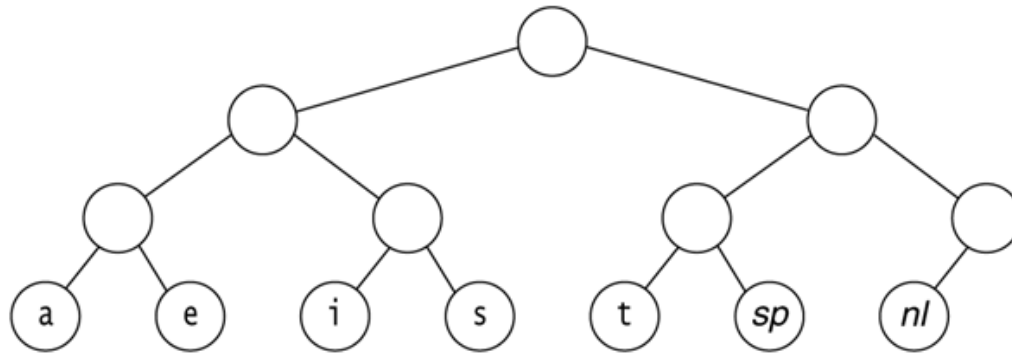


figure 12.2

Representation of the original code by a tree

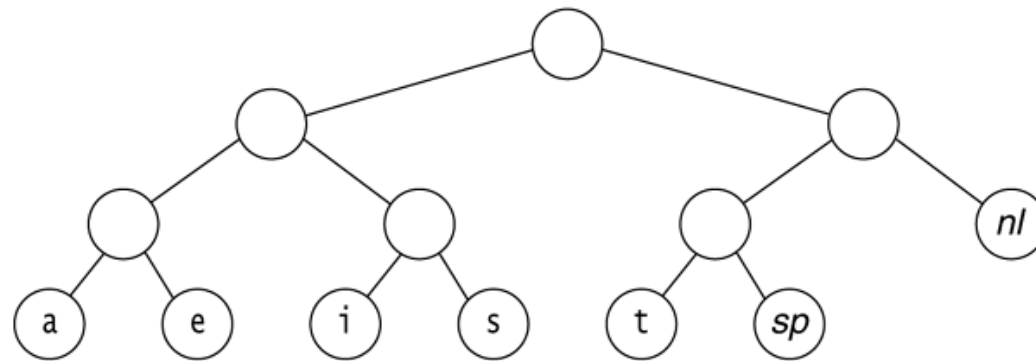


figure 12.3
A slightly better tree



figure 12.4

An optimal prefix code tree

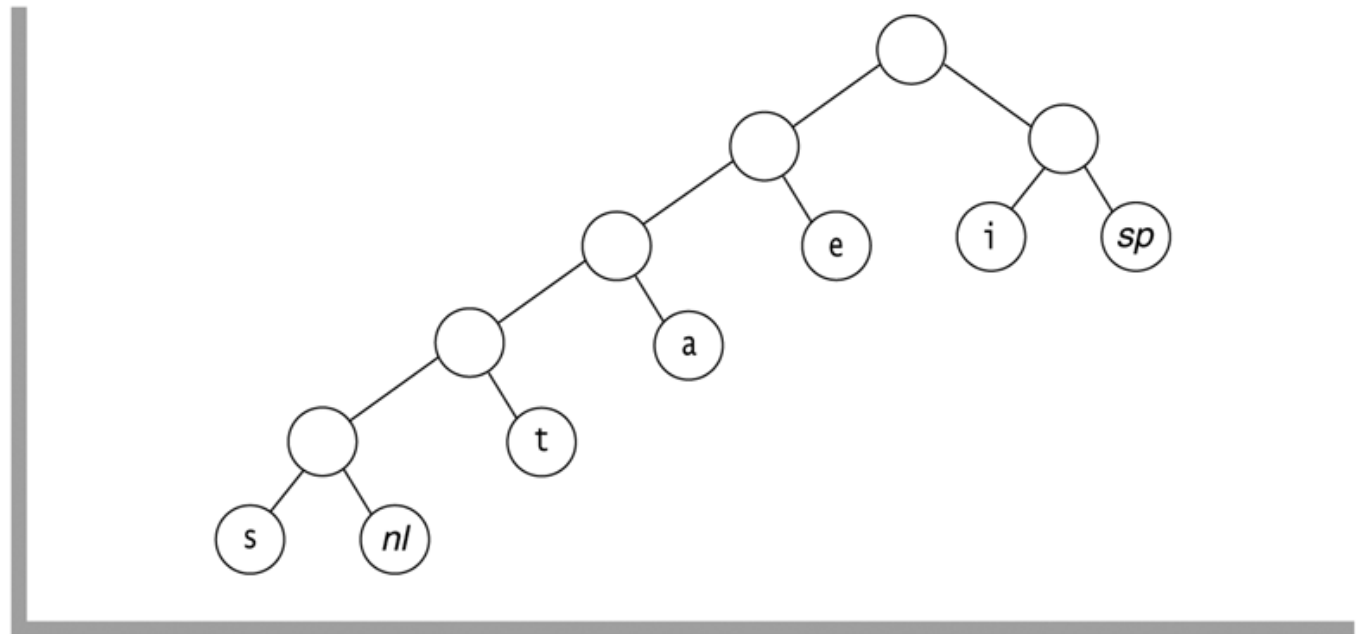




figure 12.5

Optimal prefix code

| Character | Code | Frequency | Total Bits |
|-----------|-------|-----------|------------|
| a | 001 | 10 | 30 |
| e | 01 | 15 | 30 |
| i | 10 | 12 | 24 |
| s | 00000 | 3 | 15 |
| t | 0001 | 4 | 16 |
| <i>sp</i> | 11 | 13 | 26 |
| <i>nl</i> | 00001 | 1 | 5 |
| Total | | | 146 |

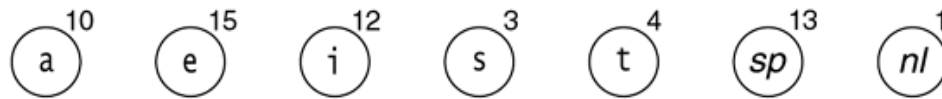


figure 12.6
Initial stage of
Huffman's algorithm

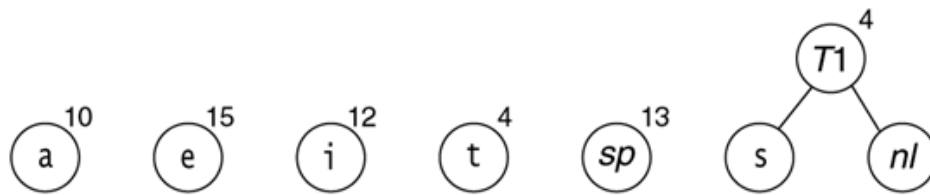


figure 12.7

Huffman's algorithm
after the first merge



figure 12.8

Huffman's algorithm
after the second
merge

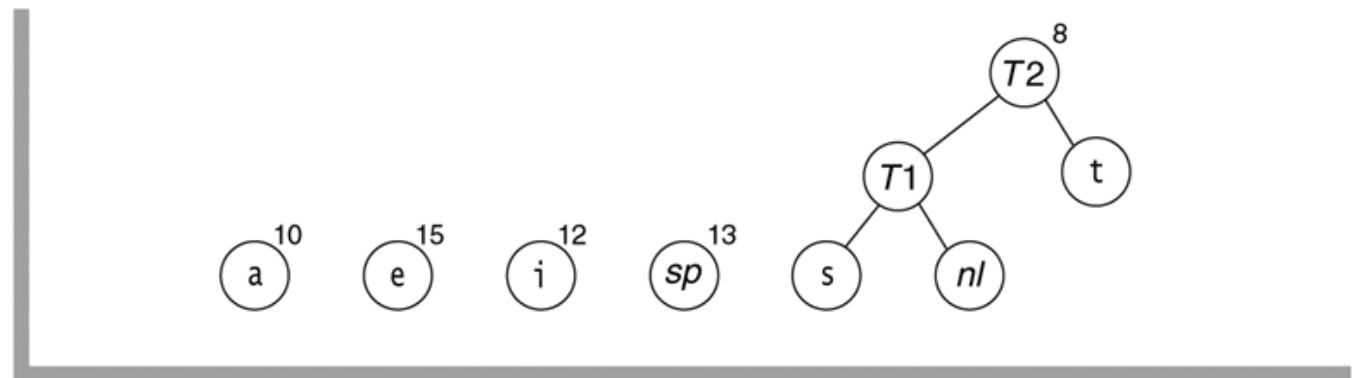




figure 12.9

Huffman's algorithm
after the third merge

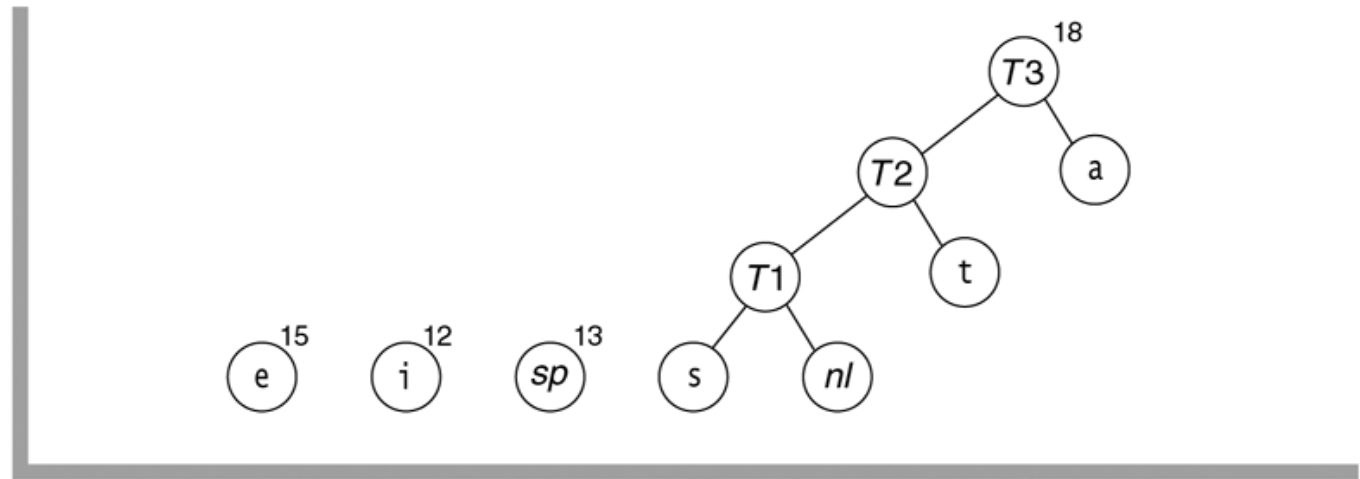




figure 12.10

Huffman's algorithm
after the fourth merge

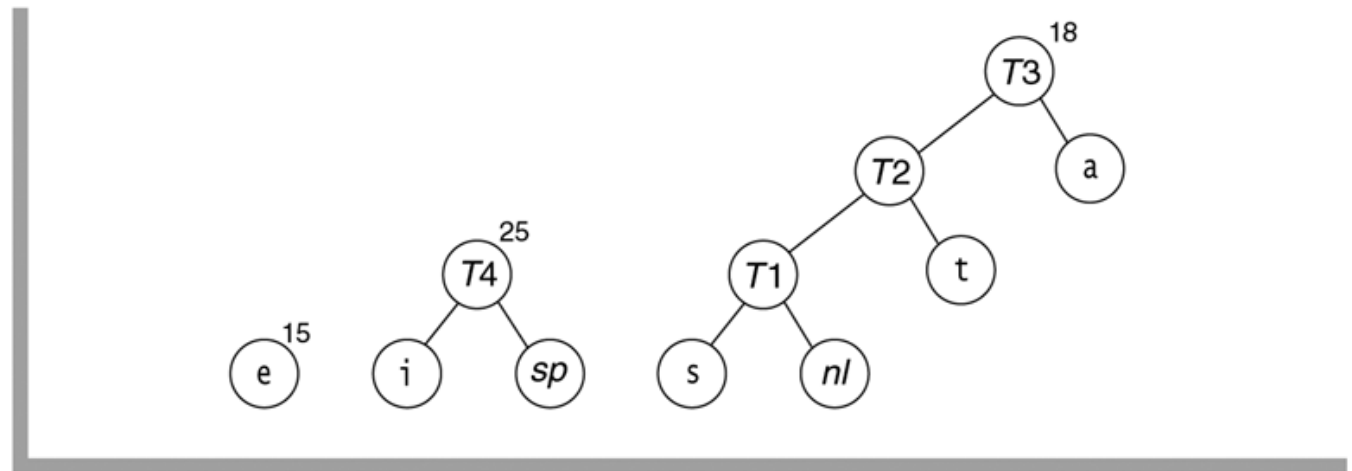
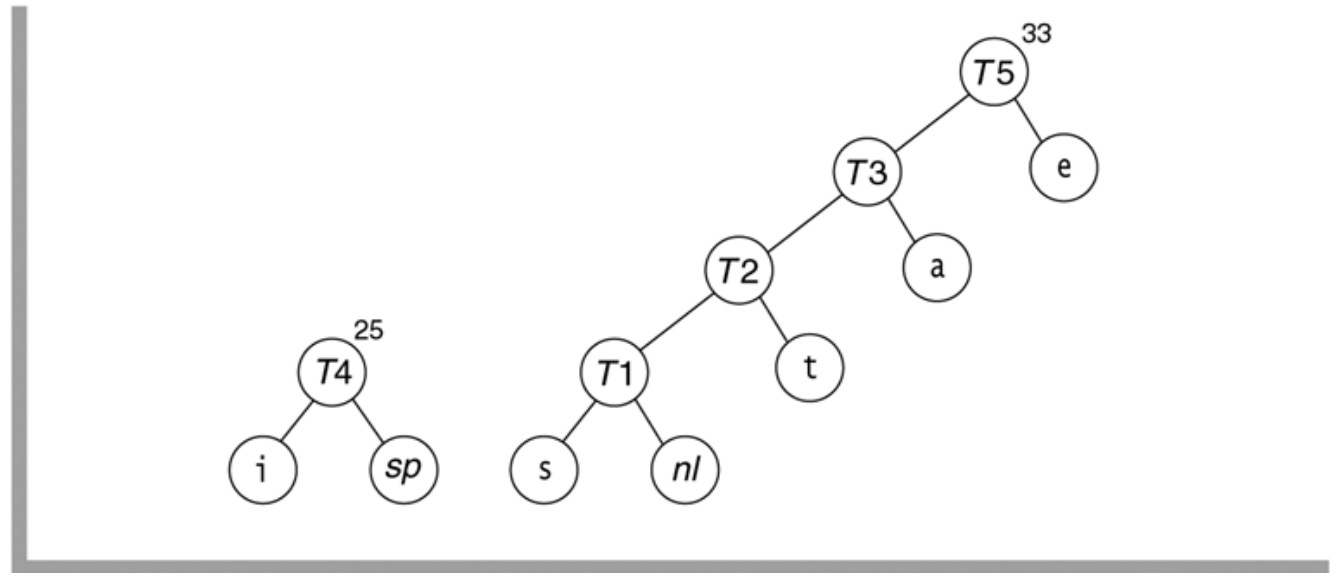




figure 12.11

Huffman's algorithm
after the fifth merge



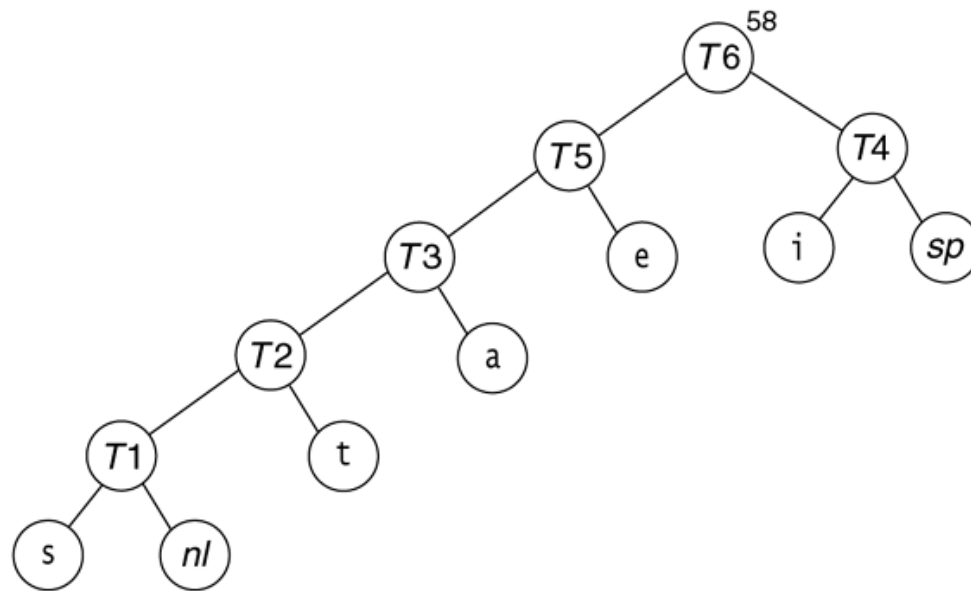


figure 12.12

Huffman's algorithm
after the final merge