

Recurrence Relations

Recurrence Relations

- Recurrence relations are used to determine the running time of recursive programs.
- Recurrence relations themselves are recursive.
- $T(0)$: time to solve problem of size 0. This is the base Case
- $T(n)$: time to solve problem of size n . This is the recursive Case

Recurrence Relation for Linear Search

- Study code
- $T(0) = 0$
- $T(n) = 1 + T(n - 1)$

Solving a Recurrence Relation through Telescoping

- $T(n) = 1 + T(n - 1)$
 $T(n - 1) = 1 + T(n - 2)$
 $T(n - 2) = 1 + T(n - 3)$
...
 $T(0) = 0$

 $T(n) = 1 + 1 + \dots + 1$ (n times)
 $T(n) = n$
 $T(n) = O(n)$

Recurrence Relation for Binary Search

- Study code
- $T(0) = 0$
- $T(n) = 2 + T(n / 2)$

Solving a Recurrence Relation through Telescoping

- $T(n) = 2 + T(n / 2)$
 $T(n / 2) = 2 + T(n / 4)$
 $T(n / 4) = 2 + T(n / 8)$
 ...
 $T(0) = 0$

$$T(n) = 2 + 2 + \dots + 2 \text{ (log n times)}$$

$$T(n) = 2 * \log(n)$$

$$T(n) = O(\log(n))$$

Recurrence Relation for Merge Sort **Best** Case

- Study code
- $T(n) = n/2 + 2T(n/2)$
 $T(n/2) = n/4 + 2T(n/4)$
 $T(n/4) = n/8 + 2T(n/8)$
...
 $T(0) = 0$
 $T(n) = n/2 + 2(n/4) + 4(n/8) + \dots$ ((log n) terms)
 $T(n) = n/2 + n/2 + n/2 + \dots$ ((log n) terms)
 $T(n) = n/2 * \log(n) = O(n \log(n))$

Recurrence Relation for Merge Sort **Worst** Case

- $T(n) = (n-1) + 2T(n/2)$
 $T(n/2) = (n/2 - 1) + 2T(n/4)$
 $T(n/4) = (n/4 - 1) + 2T(n/8)$
...
 $T(0) = 0$
 $T(n) = (n-1) + 2(n/2-1) + 4(n/4-1) + \dots$
 $= (n-1) + (n-2) + (n-4) + \dots$
 $= \log(n)*n + (-1 - 2 - 4 \dots)$ [sum of powers of 2, from 0 to log(n)-1]
 $= n*\log(n) - (n - 1)$ [$2^{(\log(n)-1)+1} - 1 = n-1$]
 $= (n-1)*\log(n) + 1$
 $= O(n \log(n))$